# Nonlinear algebra and matrix completion 

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## Motivation

## Problem

Let $\Omega \subseteq[m] \times[n]$. For a given $\Omega$-partial matrix $X \in \mathbb{C}^{\Omega}$, the low-rank matrix completion problem is

Minimize $\operatorname{rank}(M)$ subject to $M_{i j}=X_{i j}$ for all $(i, j) \in \Omega$

## Example

Let $\Omega=\{(1,1),(1,2),(2,1)\}$ and consider the following $\Omega$-partial matrix

$$
X=\left(\begin{array}{ll}
1 & 2 \\
3 & \cdot
\end{array}\right)
$$

Some applications:

- Collaborative filtering (e.g. the "Netflix problem")
- Computer vision
- Existence of MLE in Gaussian graphical models (Uhler 2012)


## State of the art: nuclear norm minimization

The nuclear norm of a matrix, denoted $\|\cdot\|_{*}$, is the sum of its singular values

## Theorem (Candès and Tao 2010)

Let $M \in \mathbb{R}^{m \times n}$ be a fixed matrix of rank $r$ that is sufficiently "incoherent." Let $\Omega \subseteq[m] \times[n]$ index a set of $k$ entries of $M$ chosen uniformly at random. Then with "high probability," $M$ is the unique solution to

$$
\begin{array}{ll}
\operatorname{minimize} & \|X\|_{*} \\
\text { subject to } & X_{i j}=M_{i j} \quad \text { for all }(i, j) \in \Omega .
\end{array}
$$

The upshot: the minimum rank completion of a partial matrix can be recovered via semidefinite programming if:

- the known entries are chosen uniformly at random
- the completed matrix is sufficiently "incoherent"

Goal: use algebraic geometry to understand the structure of low-rank matrix completion and develop methods not requiring above assumptions

## The algebraic approach

Some subsets of entries of a rank- $r$ matrix satisfy nontrivial polynomials.

## Example

If the following matrix has rank 1 , then the bold entries must satisfy the following polynomial

$$
\left(\begin{array}{lll}
\mathbf{x}_{11} & \mathbf{x}_{12} & x_{13} \\
\mathbf{x}_{21} & x_{22} & \mathbf{x}_{23} \\
x_{31} & \mathbf{x}_{32} & \mathbf{x}_{33}
\end{array}\right) \quad x_{12} x_{21} x_{33}-x_{13} x_{31} x_{11}=0
$$

Király, Theran, and Tomioka propose using these polynomials to:

- Bound rank of completion of a partial matrix from below
- Solve for missing entries


## Question

Which subsets of entries of an $m \times n$ matrix of rank $r$ satisfy nontrivial polynomials?

## Graphs and partial matrices

Subsets of entries of a matrix can be encoded by graphs:

- non-symmetric matrices $\rightarrow$ bipartite graphs
- symmetric matrices $\rightarrow$ semisimple graphs

| Mat $_{r}^{m \times n}$ | $m \times n$ matrices <br> of rank $\leq r$ | $\left(\begin{array}{ccc}5 & \cdot & \cdot \\ -4 & -2 & \cdot \\ \cdot & 8 & 3\end{array}\right)$ | r 1  <br> r 2  <br> r 3 c 1 <br> $\mathrm{c}_{2}$  |  |
| :--- | :---: | :---: | :---: | :---: |
| Sym $_{r}^{n \times n}$ | $n \times n$ symmetric <br> matrices of rank $\leq r$ | $\left(\begin{array}{ccc}7 & 4 & \cdot \\ 4 & \cdot & 9 \\ \cdot & 9 & 5\end{array}\right)$ | 0 | 3 |

- A G-partial matrix is a partial matrix whose known entries lie at the positions corresponding to the edges of $G$.
- A completion of a $G$-partial matrix $M$ is a matrix whose entries at positions corresponding to edges of $G$ agree with the entries of $M$.


## Generic completion rank

## Definition

Given a (bipartite/semisimple) graph $G$, the generic completion rank of $G$, denoted $\operatorname{gcr}(G)$, is the minimum rank of a complex completion of a $G$-partial matrix with generic entries.

| type | G | pattern | $\operatorname{gcr}(G)$ |
| :---: | :---: | :---: | :---: |
| symm | 1 ¢ $0_{2}$ | $\left(\begin{array}{cc}a_{11} & ? \\ ? & a_{22}\end{array}\right)$ | 1 |
| symm | $10 \cdot 3$ | $\left(\begin{array}{ccc}a_{11} & a_{12} & ? \\ a_{12} & a_{22} & a_{23} \\ ? & a_{23} & ?\end{array}\right)$ | 2 |
| non |  | $\left(\begin{array}{ccc}a_{11} & a_{12} & ? \\ a_{21} & ? & a_{23}\end{array}\right)$ | 1 |

## Generic completion rank

## Problem

Gain a combinatorial understanding of generic completion rank - how can one use the combinatorics of $G$ to infer $\operatorname{gcr}(G)$ ?

## Proposition (Folklore)

Given a bipartite graph $G, \operatorname{gcr}(G) \leq 1$ iff $G$ has no cycles.

## Proposition (Folklore)

Given a semisimple graph $G, \operatorname{gcr}(G) \leq 1$ iff $G$ has no even cycles, and every connected component has at most one odd cycle.


## Generic completion rank 2 - nonsymmetric case

A cycle in a directed graph is alternating if the edge directions alternate.


Alternating cycle


Non-alternating cycle

## Theorem (B.-, 2016)

Given a bipartite graph $G, \operatorname{gcr}(G) \leq 2$ if and only if there exists an acyclic orientation of $G$ that has no alternating cycle.


$$
\operatorname{gcr}(G)=2
$$


$\operatorname{gcr}(G)=3$

## Generic completion rank 2 - nonsymmetric case

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## Proof sketch

## Theorem (B.-, 2016)

Given a bipartite graph $G, \operatorname{gcr}(G) \leq 2$ if and only if there exists an acyclic orientation of $G$ that has no alternating cycle.

- Rephrase the question: describe the independent sets in the algebraic matroid underlying the variety of $m \times n$ matrices of rank at most 2
- This algebraic matroid is a restriction of the algebraic matroid underlying a Grassmannian $\operatorname{Gr}(2, N)$ of affine planes
- Algebraic matroid structure is preserved under tropicalization
- Apply Speyer and Sturmfels' result characterizing the tropicalization of $\operatorname{Gr}(2, N)$ in terms of tree metrics to reduce to an easier combinatorial problem


## Open question

Does there exist a polynomial time algorithm to check the combinatorial condition in the above theorem, or is this decision problem NP-hard?

## Issue: real vs complex

What happens when you only want to consider real completions?

## Definition

Given a bipartite or semisimple graph $G$, there may exist multiple open sets $U_{1}, \ldots, U_{k}$ in the space of real $G$-partial matrices such that the minimum rank of a completion of a partial matrix in $U_{i}$ is $r_{i}$. We call the $r_{i}$ s the typical ranks of $G$.

The graph


$$
\left(\begin{array}{cc}
a_{11} & \cdot \\
\cdot & a_{22}
\end{array}\right)
$$

In a completion to rank 1 , the missing entry $t$ must satisfy $a_{11} a_{22}-t^{2}=0$.


## Facts about typical ranks

## Proposition (B.-Blekherman-Sinn 2018)

Let $G$ be a bipartite or semisimple graph.
(1) The minimum typical rank of $G$ is $\operatorname{gcr}(G)$.
(2) The maximum typical rank of $G$ is at most $2 \operatorname{gcr}(G)$.
(3) All integers between $\operatorname{gcr}(G)$ and the maximum typical rank of $G$ are also typical ranks of $G$.

See also Bernardi, Blekherman, and Ottaviani 2015 and Blekherman and Teitler 2015.

## Case study: disjoint union of cliques

Let $K_{m} \sqcup K_{n}$ denote the disjoint union of two cliques with all loops


## Proposition (B.-Blekherman-Lee)

The generic completion rank of $K_{m} \sqcup K_{n}$ is $\max \{m, n\}$. The maximum typical rank of $K_{m} \sqcup K_{n}$ is $m+n$.

## Case study: disjoint union of cliques

## Proposition (B.-Blekherman-Lee)

The generic completion rank of $K_{m} \sqcup K_{n}$ is $\max \{m, n\}$. The maximum typical rank of $K_{m} \sqcup K_{n}$ is $m+n$.

A $\left(K_{m} \sqcup K_{n}\right)$-partial matrix looks like:

$$
M=\left(\begin{array}{cc}
A & X \\
X^{T} & B
\end{array}\right)
$$

By Schur complements:

$$
\operatorname{rank}(M)=\operatorname{rank}(A)+\operatorname{rank}\left(B-X^{\top} A^{-1} X\right)
$$

If $A \prec 0$ and $B \succ 0$, then $\operatorname{det}\left(B-X^{T} A^{-1} X\right)>0$ for real $X$.

## Corollary

Every integer between $\max \{m, n\}$ and $m+n$ is a typical rank of $K_{m} \sqcup K_{n}$.

## Case study: disjoint union of cliques

Given real symmetric matrices $A$ and $B$ of full rank, of possibly different sizes:

- $p_{A}\left(p_{B}\right)$ denotes the number of positive eigenvalues of $A(B)$
- $n_{A}\left(n_{B}\right)$ denotes the number of negative eigenvalues of $A(B)$
- the eigenvalue sign disagreement of $A$ and $B$ is defined as:

$$
\operatorname{esd}(A, B):= \begin{cases}0 & \text { if }\left(p_{A}-p_{B}\right)\left(n_{A}-n_{B}\right) \geq 0 \\ \min \left\{\left|p_{A}-p_{B}\right|,\left|n_{A}-n_{B}\right|\right\} & \text { otherwise }\end{cases}
$$

## Theorem (B.-Blekherman-Lee)

Let $M=\left(\begin{array}{cc}A & X \\ X^{T} & B\end{array}\right)$ be a generic real $K_{m} \sqcup K_{n}$-partial matrix. Then $M$ is minimally completable to rank $\max \{m, n\}+\operatorname{esd}(A, B)$.

## When full rank is typical

## Theorem (B.-Blekherman-Lee)

Let $G$ be a semisimple graph on $n$ vertices. Then $n$ is a typical rank of $G$ if and only if the complement graph of $G$ is bipartite.

If the complement is bipartite, then $n$ is a typical rank:

$$
M=\left(\begin{array}{cc}
A & X \\
X^{T} & B
\end{array}\right)
$$

By Schur complements:

$$
\operatorname{rank}(M)=\operatorname{rank}(A)+\operatorname{rank}\left(B-X^{\top} A^{-1} X\right)
$$

so if $A \prec 0$ and $B \succ 0$, then $\operatorname{det}\left(B-X^{T} A^{-1} X\right)$ is strictly positive.

## When full rank is typical

## Theorem (B.-Blekherman-Lee)

Let $G$ be a semisimple graph on $n$ vertices. Then $n$ is a typical rank of $G$ if and only if the complement graph of $G$ is bipartite.

If complement is not bipartite, then $n$ is not a typical rank:

- A graph is bipartite if and only if it is free of odd cycles
- If complement graph is an odd cycle, then determinant of a G-partial matrix, viewed as a polynomial in the unknown entires, has odd degree
- Deleting edges from a graph will not increase maximum typical rank.

$$
\left(\begin{array}{ccccc}
a_{11} & \mathbf{x} & a_{13} & a_{14} & \mathbf{t} \\
\mathbf{x} & a_{22} & \mathbf{y} & a_{24} & a_{25} \\
a_{13} & \mathbf{y} & a_{33} & \mathbf{z} & a_{35} \\
a_{14} & a_{24} & \mathbf{z} & a_{44} & \mathbf{w} \\
\mathbf{t} & a_{25} & a_{35} & \mathbf{w} & a_{55}
\end{array}\right)
$$

## Typical ranks for nonsymmetric matrices: some examples

The following bipartite graph has 2 and 3 as typical ranks.


$$
\left(\begin{array}{cccc}
? & a_{12} & a_{13} & a_{14} \\
a_{21} & ? & a_{23} & a_{24} \\
a_{31} & a_{32} & ? & a_{34} \\
a_{41} & a_{42} & a_{43} & ?
\end{array}\right)
$$

Let $\operatorname{mtr}(G)$ denote the maximum typical rank of $G$.

## Theorem (B.-Blekherman-Sinn)

Let $G$ be obtained by gluing two bipartite graphs $G_{1}$ and $G_{2}$ along a complete bipartite subgraph $K_{m, n}$. If

$$
\max \left\{\operatorname{mtr}\left(G_{1}\right), \operatorname{mtr}\left(G_{2}\right)\right\} \geq \max \{m, n\},
$$

then $\operatorname{mtr}(G)=\max \left\{\operatorname{mtr}\left(G_{1}\right), \operatorname{mtr}\left(G_{2}\right)\right\}$. The same is true for generic completion rank.

## Open question

Does there exist a bipartite graph that has more than two typical ranks?

## Empty $k$-cores

The $k$-core of a graph $G$ is the graph obtained by iteratively removing vertices of degree $k-1$ or less. The 2-core of the graph below is empty.


Theorem (B.-, Blekherman, Sinn)
Let $G$ be bipartite. If the $k$-core of $G$ is empty, then all typical ranks of $G$ are at most $k-1$.

## Corollary

Let $G$ be bipartite. Then the maximum typical rank of $G$ is $2 \operatorname{gcr}(G)-1$.

## Open question

Which bipartite graphs of generic completion rank 2 also have 3 as a typical rank?

## Conclusion

- All generic $G$-partial matrices can be completed to rank $\operatorname{gcr}(G)$ over $\mathbb{C}$
- We can characterize all the bipartite graphs with generic completion rank $\leq 2$ (semisimple case is still open)
- Over the reals, a graph can have many typical ranks

Open problems:

- Find a polynomial-time algorithm to decide if a given bipartite graph has an acyclic orientation with no alternating cycle, or prove that this decision problem is NP-hard
- Find a bipartite graph that exhibits three or more typical ranks
- Characterize the graphs with generic completion rank 2 that also exhibit 3 as a typical rank


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